

## Cycles of Even Length in Graphs

J. A. BONDY

*University of Waterloo, Waterloo, Ontario, Canada*

AND

M. SIMONOVITS

*Eötvös Loránd University, Budapest, Hungary*

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In this paper we solve a conjecture of P. Erdős by showing that if a graph  $G^n$  has  $n$  vertices and at least  $100kn^{1+1/k}$  edges, then  $G$  contains a cycle  $C^{2l}$  of length  $2l$  for every integer  $l \in [k, kn^{1/k}]$ . Apart from the value of the constant this result is best possible. It is obtained from a more general theorem which also yields corresponding results in the case where  $G^n$  has only  $cn(\log n)^\alpha$  edges ( $\alpha > 1$ ).

### 0. NOTATION

The graphs considered in this paper are finite and have neither loops nor multiple edges. The number of edges of a graph  $G$  will be denoted by  $e(G)$ . The number of vertices will be either denoted by  $v(G)$  or indicated by a superscript; thus  $G^n$  is always a graph on  $n$  vertices.  $C^k$  denotes the cycle of length  $k$ .

### 1. INTRODUCTION

P. Erdős, in [4], published without proof the following

**THEOREM.** *There exists a  $c_k$  and an  $n_0(k)$  such that, if*

$$e(G^n) > c_k n^{1+1/k} \quad \text{and} \quad n > n_0, \quad (1)$$

*then*

$$C^{2k} \subset G^n.$$

Later, Erdős asked whether (1) implies

$$C^{2l} \subset G^n \text{ for every integer } l \in [k, n^{1/k}].$$

(In [5] he proved a weaker form of this conjecture for  $k = 2$ ). We shall prove

THEOREM 1. *If*

$$e(G^n) > 100kn^{1+1/k}, \quad (2)$$

*then*

$$C^{2l} \subset G^n \text{ for every integer } l \in [k, kn^{1/k}].$$

*Remark 1.* It is reasonable to conjecture the existence of a function  $f$  such that, for all sufficiently large  $n$ , there is a graph  $S^n$  with  $[f(k)n^{1+1/k}]$  edges that does not contain a  $C^{2k}$ ; this is known to be the case for  $k = 2, 3$ , and 5 ([3], [7], [1], [8]). Therefore (at least for these values of  $k$ ), condition (2) cannot be replaced by

$$e(G^n) > f(k)n^{1+1/k}.$$

In this sense our theorem is sharp.

On the other hand, if  $Z^n$  is the union of approximately  $(1/k)n^{1+1/k}$  complete graphs on  $[kn^{1/k}]$  vertices, then

$$e(Z^n) \approx kn^{1+1/k};$$

but  $Z^n$  contains no cycle of length greater than  $kn^{1/k}$ . Therefore, if  $e(G) \approx kn^{1+1/k}$ , the existence of a  $C^{2l}$  in  $G^n$  for  $l = [kn^{1/k}]$  cannot be ensured, and this again shows the sharpness of our theorem.

*Remark 2.* In particular (for the case  $k = 2$ ), Theorem 1 tells us that the order of magnitude of  $e(G^n)$  which forces  $G^n$  to contain a  $C^4$  also forces  $G^n$  to contain all the even cycles  $C^{2l}$ ,  $l = 2, 3, \dots, 2n^{1/2}$ . A similar phenomenon is established in a paper of J. A. Bondy [2], where it is shown that, if  $G^n$  has enough edges to force a triangle (that is, if  $e(G^n) > (n^2/4)$ ), then  $G^n$  must contain all cycles  $C^l$ ,  $l = 3, 4, \dots, [(n+3)/2]$ .

Theorem 1 is an easy consequence of a slightly more general theorem.

THEOREM 1\*. *Let  $E = e(G^n)$ . Then  $C^{2l} \subset G^n$  for every integer  $l \geq 2$  satisfying*

$$l \leq \frac{E}{100n}, \quad ln^{1/l} \leq \frac{E}{10n}.$$

Besides Theorem 1, another consequence of Theorem 1\* is

THEOREM 2. *There exists a function  $g$  such that, if*

$$e(G^n) \geq g(\epsilon) n(\log n)^{1+\epsilon},$$

*then*

$$C^{2l} \subset G^n \quad \text{for every integer } l \in \left[ \frac{\log n}{\epsilon \log \log n}, (\log n)^{1+\epsilon} \right].$$

## 2. BASIC LEMMAS

A coloring (not necessarily proper) of the vertices of a graph  $G$  is  $t$ -periodic if the end-vertices of any (simple) path of length  $t$  in  $G$  have the same color.

LEMMA 1. Let  $t$  be a positive integer, and let  $G$  be a connected graph for which

$$e(G) \geq 2tv(G). \quad (3)$$

*Then the number of colors in any  $t$ -periodic coloring of  $G$  is at most two.*

*Proof.* (i) First we show that any graph  $G$  with  $e(G) \geq 2tv(G)$  contains two adjacent vertices joined by two vertex-disjoint paths, each of length at least  $t$ . The technique we use is due to Pósa.

In the case where each vertex has valence at least  $2t$ , we can find such a  $\theta$ -graph in the following way. Let a longest path in  $G$  be  $x_1 \dots x_m$ . Then  $x_1$  is adjacent only to vertices of this path, say to  $x_{i_1}, x_{i_2}, \dots, x_{i_r}$ , where

$$2 = i_1 < i_2 < \dots < i_r \text{ and } r \geq 2t.$$

The path  $x_1x_2 \dots x_{i_{2t}}$  together with the edges  $x_1x_{i_1}$  and  $x_1x_{i_{2t}}$  form the desired  $\theta$ -graph.

The general case, when there may be vertices of valence less than  $2t$ , can now be proved by induction on  $v(G)$ . For  $0 < v(G) \leq 4t$  it is trivial that (3) cannot be satisfied, and so there is nothing to prove here. If  $v(G) = 4t + 1$ ,  $G$  must be complete and clearly contains a  $\theta$ -graph of the desired type. Suppose now that every graph that satisfies (3) and has  $k \geq 4t + 1$  vertices contains such a  $\theta$ -graph, and let  $G$  be a graph on  $k + 1$  vertices with some vertex  $x$  of valence less than  $2t$ . Then

$$e(G - x) > e(G) - 2t \geq 2tv(G) - 2t = 2tv(G - x).$$

Thus, by the induction hypothesis,  $G - x$  contains a  $\theta$ -graph of the desired type and hence so also does  $G$ .

(ii) Let the three cycles of such a  $\theta$ -graph be  $C_1, C_2, C_3$  with lengths  $l_1, l_2, l_3$ , respectively. Clearly, the restrictions of the  $t$ -periodic coloring of  $G$  to the  $\theta$ -graph and to each cycle  $C_i$  are also  $t$ -periodic. Let  $t_i$  be the smallest integer such that  $C_i$  is  $t_i$ -periodic,  $i = 1, 2, 3$ . It is easy to see that any period on one cycle induces the same period on the other cycles and therefore

$$t_1 = t_2 = t_3;$$

also,  $t_i \mid l_i, i = 1, 2, 3$ . If  $C_3$  is the longest of the three cycles, then

$$l_1 \div l_2 - l_3 = 2.$$

Setting  $t_i = t^*, i = 1, 2, 3$ , we find that  $t^* \mid 2$  and hence that  $t^* = 1$  or  $t^* = 2$ . Therefore, the number of colors in the  $\theta$ -graph is at most two.

(iii) Because  $G$  is connected, each vertex of  $G$  is joined to some vertex of this  $\theta$ -graph by a path of length  $kt$ , for some integer  $k$ , and hence has the same color as this vertex. It follows that the number of colors in the whole graph  $G$  is also at most two. This completes the proof of the lemma.

It is, in fact, easy to show that either  $G$  is bipartite with the natural coloring (trivially a 2-periodic coloring), or else  $G$  is unicolored.

LEMMA 2. *Let  $G^n$  be a bipartite graph in which every vertex has valence at least  $s = \max\{5ln^{1/l}, 50l\}$ . Then  $G^n$  contains a  $C^{2l}$ .*

*Proof.* Choose an arbitrary vertex  $x$  of  $G^n$  and let  $V_i$  be the set of vertices at distance  $i$  from  $x$ . Since  $G^n$  is bipartite, each set  $V_i$  is an independent set.

Suppose that  $G^n$  contains no  $C^{2l}$ . We shall show that this implies that, for  $1 \leq i \leq l$ ,

$$\frac{|V_i|}{|V_{i-1}|} \geq \frac{s}{5l}, \quad (4)$$

thus leading to the contradiction that  $v(G^n) > n$ , (since  $s \geq 5ln^{1/l}$  and, consequently,  $|V_i| \geq n^{1/l} |V_{i-1}|$ ).

We prove (4) by induction on  $i$ . It is trivial for  $i = 1$  since the vertex  $x$  has valence at least  $s$ . Suppose that it is true for  $i - 1$ . Let  $H_1, H_2, \dots, H_q$  be the components of the subgraph  $H$  of  $G^n$  induced by  $V_{i-1} \cup V_i$ , and let  $W_j$  be the set of vertices of  $H_j$  that are on level  $i - 1$ , that is, in  $V_{i-1}$  (see Figure I).

A path  $x_1x_2 \dots x_m$  in  $G^n$  will be called *monotonic* if the distance between  $x$  and  $x_i$  is monotonic. (This means that a monotonic path crosses any level at most once.)

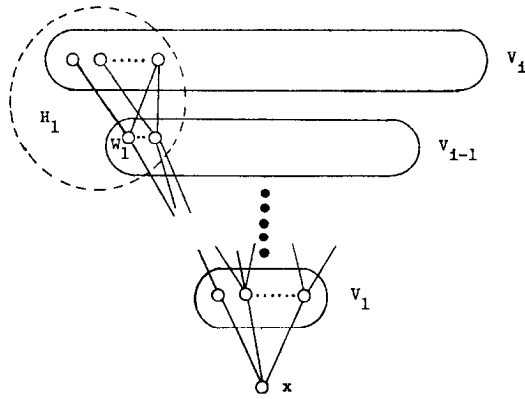


FIGURE 1

We shall show that  $e(H_1) < 4lv(H_1)$ . This is trivial if  $W_1$  has just one vertex, so assume that  $W_1$  has at least two vertices. Let  $a \in V_p$  be a vertex of  $G^n$  such that

- (i) there are two monotonic paths  $P_1, P_2$  joining  $a$  to  $W_1$  which have just the vertex  $a$  in common,
- (ii)  $p$  is the minimum subject to (i).

First we show that each vertex of  $W_1$  is joined to  $a$  by a monotonic path. For  $y \in W_1$  is joined to  $x$  by a monotonic path  $P_3$  and, by the minimality of  $p$ ,  $P_3$  must intersect  $P_1$  in some vertex  $z$ . The path consisting of the section of  $P_3$  between  $y$  and  $z$  and the section of  $P_1$  between  $z$  and  $a$  is a monotonic path from  $y$  to  $a$ . This is illustrated in Figure 2.

We now assign colors red and blue to the vertices of  $W_1$  in such a way

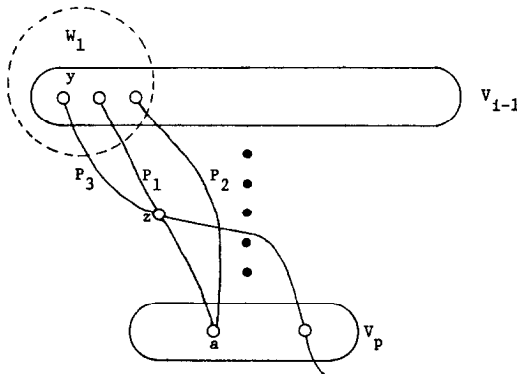


FIGURE 2

that, if two vertices have different colors, then they are joined to  $a$  by vertex-disjoint monotonic paths. This is done as follows. Each vertex of  $W_1$  that can be joined to  $a$  by a monotonic path disjoint from  $P_2$  is colored red; all other vertices of  $W_1$  are colored blue. To see that this coloring has the required property, let  $x_1$  and  $x_2$  be vertices of  $W_1$  colored red and blue, respectively, let  $P_1'$  be a monotonic path from  $x_1$  to  $a$  disjoint from  $P_2$ , and let  $P_2'$  be a monotonic path from  $x_2$  to  $a$ . Moving along  $P_2'$  from  $x_2$  towards  $a$ , let  $v$  be the first vertex of  $(P_1' \cup P_2) - a$  encountered (see Figure 3). Because  $x_2$  has the color blue, such a  $v$  exists;  $v$  cannot belong

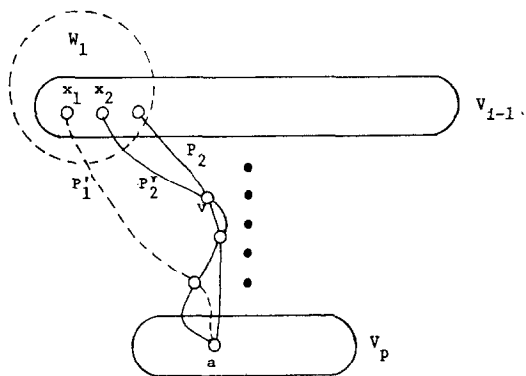


FIGURE 3

to  $P_1'$  for then the section of  $P_2'$  between  $x_2$  and  $v$  together with the section of  $P_1'$  between  $v$  and  $a$  would constitute a monotonic path from  $x_2$  to  $a$  disjoint from  $P_2$ , contradicting the assumption that  $x_2$  is colored blue. But then  $v \in P_2$  and we have a monotonic path  $x_2 P_2' v P_2 a$  disjoint from  $P_1'$ .

We now color the vertices of  $H_1$  in  $V_i$  green and show that this coloring of  $H_1$  is  $t$ -periodic with  $t = 2(l - i + p \pm 1)$ . For, since  $t$  is even, if one end-vertex of a path of length  $t$  in  $H_1$  is green, then so is the other. Also, there can be no path of length  $t$  joining a red and a blue vertex, because, if a red  $x_1$  were joined to a blue  $x_2$  by such a path, this path together with vertex-disjoint monotonic paths from  $x_1$  to  $a$  and from  $x_2$  to  $a$  would form a cycle of length  $2l$ . Therefore, the coloring of  $H_1$  is indeed  $t$ -periodic. Since three colors are used in this coloring, Lemma 1 implies that

$$e(H_1) < 2tv(H_1) < 4lv(H_1).$$

Arguing similarly for  $H_2, \dots, H_q$ , we obtain

$$e(H_j) < 4lv(H_j), \quad j = 1, \dots, q,$$

and, since the  $H_j$  are the components of  $H$ ,

$$e(H) < 4lv(H). \quad (5)$$

Let  $H^*$  denote the subgraph of  $G^n$  induced by  $V_{i-2} \cup V_{i-1}$ . Then, similarly,

$$e(H^*) < 4lv(H^*), \quad (6)$$

and, by the induction hypothesis,

$$\frac{|V_{i-1}|}{|V_{i-2}|} \geq \frac{s}{5l}. \quad (7)$$

But, clearly, since each vertex of  $G^n$  has valence at least  $s$ ,

$$e(H) + e(H^*) \geq s |V_{i-1}|.$$

Therefore, by (5) and (6),

$$\begin{aligned} 4l(|V_{i-1}| + |V_i| + |V_{i-2}| + |V_{i-1}|) \\ = 4l(v(H) + v(H^*)) > e(H) + e(H^*) \geq s |V_{i-1}| \end{aligned}$$

and so

$$|V_i| > \frac{1}{4l} ((s - 8l) |V_{i-1}| - 4l |V_{i-2}|).$$

Using (7) we obtain

$$|V_i| > \frac{1}{4l} \left( s - 8l - \frac{20l^2}{s} \right) |V_{i-1}|$$

and, therefore, since  $s \geq 50l$ ,

$$\frac{|V_i|}{|V_{i-1}|} > \frac{1}{4l} (s - 9l) > \frac{1}{4l} \cdot \frac{4s}{5} = \frac{s}{5l},$$

as desired.

### 3. MAIN THEOREM

We are now in a position to prove Theorem 1\*. First we recall its statement.

**THEOREM 1\*.** *Let  $E = e(G^n)$ . Then  $C^{2l} \subset G^n$  for every integer  $l \geq 2$  satisfying*

$$l \leq \frac{E}{100n}, \quad ln^{1/l} \leq \frac{E}{10n}. \quad (8)$$

*Proof* (by induction on  $n$ ). For  $n = 1$  the theorem is trivial, since condition (8) cannot be satisfied in this case. We now suppose that the theorem has been proved for all graphs on  $n - 1$  vertices. Let  $G^n$  be a graph on  $n$  vertices and let  $l \geq 2$  be an integer satisfying (8).

It has been shown by Erdős [6] that any graph  $G$  contains a bipartite spanning subgraph  $H$  with  $e(H) \geq e(G)/2$ ; in fact  $H$  can be chosen so that each vertex has valence in  $H$  at least half its valence in  $G$ .

So let  $H^n$  be such a bipartite spanning subgraph of  $G^n$ . If each vertex of  $H^n$  has valence at least  $E/2n$  then, by Lemma 2, we have that, for every integer  $l$  such that

$$\max\{5ln^{1/l}, 50l\} \leq E/2n,$$

$H^n$  contains a cycle of length  $2l$ . Thus, in this case, Theorem 1\* is proved.

So suppose now that some vertex  $w$  of  $H^n$  has valence less than  $E/2n$ . By the choice of  $H^n$ ,  $w$  has valence less than  $E/n$  in  $G^n$ . Let  $G^{n-1} = G^n - w$ . By the induction hypothesis,  $G^{n-1}$  contains a cycle of length  $2l$  for every integer  $l$  satisfying

$$l \leq \frac{e(G^{n-1})}{100(n-1)}, \quad l(n-1)^{1/l} \leq \frac{e(G^{n-1})}{10(n-1)}.$$

But if  $l$  satisfies (8) with  $G^n$ , then it also satisfies (8) with  $G^{n-1}$  since,

(a) if  $l \leq e(G^n)/100n$ , then

$$l \leq \frac{e(G^n)}{100n} = \frac{e(G^n) - e(G^n)/n}{100(n-1)} \leq \frac{e(G^{n-1})}{100(n-1)}$$

(since  $w$  has valence less than  $e(G^n)/n$ ).

(b) if  $ln^{1/l} \leq e(G^n)/10n$ , then

$$l(n-1)^{1/l} \leq ln^{1/l} \leq \frac{e(G^n)}{10n} = \frac{e(G^n) - e(G^n)/n}{10(n-1)} \leq \frac{e(G^{n-1})}{10(n-1)}.$$

Hence  $G^{n-1}$ , and therefore also  $G^n$ , contains a cycle of length  $2l$  for every integer  $l$  satisfying (8). This completes the proof.

Perhaps, by other methods, Theorem 1\* could be improved so as to be meaningful for

$$E \geq \frac{cn \log n}{\log \log n}.$$



However, this would then be the best possible result since, if  $c^*$  is small, there exists no fixed  $l$  such that every graph on  $n$  vertices and with  $(c^*n \log n)/(\log \log n)$  edges has a cycle of length  $2l$ .

*Remark 3.* One can find an  $l$  satisfying (8) in Theorem 1\* if and only if

$$E \geq \frac{100n \log n}{\log 10}. \quad (9)$$

If (9) holds, then (8) is satisfied for all values of  $l$  in an interval. The upper end of this interval is  $E/100n$ . The lower end can be determined in the following way:

For a fixed  $n$  the function  $y = xn^{1/x}$  is strictly decreasing in  $(0, \log n]$ . Let  $\phi_n(y) = x$  denote its inverse. Then  $\phi_n(E/10n)$  is the lower end of our interval.  $\phi_n(E/10n)$  is a transcendental function but one can easily give good approximations for it using the iteration

$$\psi_{n,1}(y) = \frac{\log n}{\log y}, \quad \psi_{n,k}(y) = \frac{\log n}{\log y - \log \psi_{n,k-1}}.$$

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